



Enumerating the edge-colourings and total colourings of a regular graph

Stéphane Bessy, Frédéric Havet

► To cite this version:

Stéphane Bessy, Frédéric Havet. Enumerating the edge-colourings and total colourings of a regular graph. [Research Report] RR-7652, INRIA. 2011. inria-00602188

HAL Id: inria-00602188

<https://hal.inria.fr/inria-00602188>

Submitted on 21 Jun 2011

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.



INSTITUT NATIONAL DE RECHERCHE EN INFORMATIQUE ET EN AUTOMATIQUE

Enumerating the edge-colourings and total colourings of a regular graph

S. Bessy and F. Havet

N° 7652

June 2011

Thème COM

*Rapport
de recherche*

Enumerating the edge-colourings and total colourings of a regular graph*

S. Bessy[†] and F. Havet[‡]

Thème COM — Systèmes communicants
Équipe-Projet Mascotte

Rapport de recherche n° 7652 — June 2011 — 14 pages

Abstract: In this paper, we are interested in computing the number of edge colourings and total colourings of a graph. We prove that the maximum number of k -edge-colourings of a k -regular graph on n vertices is $k \cdot (k-1)^{n/2}$. Our proof is constructive and leads to a branching algorithm enumerating all the k -edge-colourings of a k -regular graph using a time $O^*((k-1)^{n/2})$ and polynomial space. In particular, we obtain a algorithm on time $O^*(2^{n/2}) = O^*(1.4143^n)$ and polynomial space to enumerate all the 3-edge-colourings of a cubic graph, improving the running time of $O^*(1.5423^n)$ of the algorithm due to Golovach et al. [9]. We also show that the number of 4-total-colourings of a connected cubic graph is at most $3.2^{3n/2}$. Again, our proof yields a branching algorithm to enumerate all the 4-total-colourings of a connected cubic graph.

Key-words: edge colouring, total colouring, enumeration, (s,t)-ordering, regular graph

* This work is supported by the French *Agence Nationale de la Recherche* under reference AGAPE ANR-09-BLAN-0159.

[†] Université Montpellier 2 - CNRS, LIRMM. e-mail: Stephane.Bessy@lirmm.fr.

[‡] Projet Mascotte, I3S (CNRS, UNSA) and INRIA, Sophia Antipolis.
email: Frederic.Havet@inria.fr.

Enumeration des arête-colorations et colorations totales d'un graphe régulier

Résumé : Dans ce rapport, nous nous intéressons au calcul du nombre d'arête-colorations et du nombre de colorations totales d'un graphe. Nous montrons que le nombre maximum de k -arête-colorations d'un graphe k -régulier à n sommets est $k \cdot (k-1)!^{n/2}$. Notre preuve est constructive et donne un algorithme de branchement pour énumérer toutes les k -arête-colorations d'un graphe k -régulier en temps $O^*((k-1)!^{n/2})$ et espace polynomial. En particulier, nous obtenons un algorithme en temps $O^*(2^{n/2}) = O^*(1.4143^n)$ et espace polynomial qui énumère toutes les 3-arête-colorations d'un graphe cubique, améliorant ainsi le temps $O^*(1.5423^n)$ de l'algorithme de Golovach et al. [9]. Nous montrons également que le nombre de 4-colorations totales d'un graphe cubique connexe est au plus $3 \cdot 2^{3n/2}$. Là aussi, la preuve donne un algorithme de branchement pour énumérer toutes les 4-colorations totales d'un graphe cubique connexe.

Mots-clés : arête-coloration, coloration totale, énumération, (s,t)-ordre, graphe régulier

1 Introduction

We refer to [4] for standard notation and concepts. In this paper, all the considered graphs are loopless, but may have parallel edges. A graph with no parallel edges is said to be *simple*. Let G be a graph. We denote by $n(G)$ the number of vertices of G , and for all integer k , we denote by $n_k(G)$ the number of degree k vertices of G . Often, when the graph G is clearly understood, we abbreviate $n(G)$ to n and $n_k(G)$ to n_k .

Graph colouring is one of the classical subjects in graph theory. See for example the book of Jensen and Toft [11]. From an algorithmic point of view, for many colouring type problems, like vertex-colouring, edge colouring and total colouring, the existence problem asking whether an input graph has a colouring with an input number of colours is NP-complete. Even more, these colouring problems remain NP-complete when the question is whether there is a colouring of the input graph with a fixed (and greater than 2) number of colours [8, 10, 16].

Exact algorithms to solve NP-hard problems are a challenging research subject in graph algorithms. Many papers on exact exponential time algorithms have been published in the last decade. One of the major results is the $O^*(2^n)$ inclusion-exclusion algorithm to compute the chromatic number of a graph found independently by Björklund, Husfeldt [2] and Koivisto [13]. This approach may also be used to establish a $O^*(2^n)$ algorithm to count the k -colourings and to compute the chromatic polynomial of a graph. It also implies a $O^*(2^m)$ algorithm to count the k -edge-colourings and a $O^*(2^{n+m})$ algorithm to count the k -total-colourings of a given graph.

Since edge colouring and total colouring are particular cases of vertex-colouring, a natural question is to ask if faster algorithms than the general one may be designed in these cases. For instance, very recently Björklund et al. [3] showed how to detect whether a k -regular graph admits a k -edge-colouring in time $O^*(2^{(k-1)n/2})$.

The existence problem asking whether a graph has a colouring with a fixed and small number k of colours also attracted a lot of attention. For vertex-colourability the fastest algorithm for $k = 3$ has running time $O^*(1.3289^n)$ and was proposed by Beigel and Eppstein [1], and the fastest algorithm for $k = 4$ has running time $O^*(1.7272^n)$ and was given by Fomin et al. [7]. They also established algorithms for counting k -vertex-colourings for $k = 3$ and 4. The existence problem for a 3-edge-colouring is considered in [1, 14, 9]. Kowalik [14] gave an algorithm deciding if a graph is 3-edge-colourable in time $O^*(1.344^n)$ and polynomial space and Golovach et al. [9] presented an algorithm counting the number of 3-edge-colourings of a graph in time $O^*(3^{n/6}) = O^*(1.201^n)$ and exponential space. Golovach et al. [9] also showed a branching algorithm to enumerate all the 3-edge-colourings of a connected cubic graph of running time $O^*(25^{n/8}) = O^*(1.5423^n)$ using polynomial space. In particular, this implies that every connected cubic graph of order n has at most $O(1.5423^n)$ 3-edge-colourings. They give an example of a connected cubic graph of order n having $\Omega(1.2820^n)$ 3-edge-colourings. In Section 2, we prove that a connected cubic graph of order n has at most $3 \cdot 2^{n/2}$ 3-edge-colourings and give an example

reaching this bound. Our proof can be translated into a branching algorithm to enumerate all the 3-edge-colourings of a connected cubic graph of running time $O^*(2^{n/2}) = O^*(1.4143^n)$ using polynomial space. Furthermore, we extend our result proving that every k -regular connected graph of order n admits at most $k \cdot (k-1!)^{n/2}$ k -edge-colourings. And, similarly, we derive a branching algorithm to enumerate all the k -edge-colourings of a connected k -regular graph of running time $O^*((k-1!)^{n/2})$ using polynomial space.

Regarding total colouring, very little has been done. Golovach et al. [9] showed a branching algorithm to enumerate the 4-total-colourings of a connected cubic graph in time $O^*(2^{13n/8}) = O^*(3.0845^n)$, implying that the maximum number of 4-total-colourings in a connected cubic graph of order n is at most $O^*(2^{13n/8}) = O^*(3.0845^n)$. In Section 3, we lower this bound to $3 \cdot 2^{3n/2} = O(2.8285^n)$. Again, our proof yields a branching algorithm to enumerate all the 4-total-colourings of a connected cubic graph of running time $O^*(2.8285^n)$ using polynomial space.

2 Edge colouring

A (*proper*) *edge colouring* of a graph is a colouring of its edges such that two adjacent edges receive different colours. An edge colouring with k colours is a k -*edge-colouring*. We denote by $c_k(G)$ the number of k -edge-colourings of a graph G .

2.1 General bounds for k -regular graphs

In this section, we are interested in computing the number of k -edge-colourings of k -regular graphs. We start by computing exactly the number of 3-edge-colourings of the cycles.

Proposition 1. *Let C_n be the cycle of length n .*

$$c_3(C_n) = \begin{cases} 2^n + 2, & \text{if } n \text{ is even,} \\ 2^n - 2, & \text{if } n \text{ is odd.} \end{cases}$$

Proof. By induction on n . It is easy to check that $c_3(C_2) = c_3(C_3) = 6$.

Let $C_n = (v_1, v_2, \dots, v_n)$. Let A be the set of 3-edge-colourings of C_n such that $c(v_{n-1}v_n) \neq c(v_1v_2)$ and B the set of 3-edge-colourings of C_n such that $c(v_{n-1}v_n) = c(v_1v_2)$. The 3-edge-colourings of A are in one-to-one correspondence with those of C_{n-1} and the pair of colourings of B agreeing everywhere except on v_nv_1 are in one-to-one correspondence with the 3-edge-colourings of C_{n-2} . Thus $c_3(C_n) = c_3(C_{n-1}) + 2c_3(C_{n-2})$. Hence, if n is even, then $c_3(C_n) = 2^{n-1} - 2 + 2(2^{n-2} + 2) = 2^n + 2$, and if n is odd, then $c_3(C_n) = 2^{n-1} + 2 + 2(2^{n-2} - 2) = 2^n - 2$. \square

Let us now present our method which is based on a classical tool: (s, t) -ordering.

Definition 2. Let G be a graph and s and t be two distinct vertices of G . An (s, t) -ordering of G is an ordering of its vertices v_1, \dots, v_n such that $s = v_1$ and $t = v_n$, and for all $1 < i < n$, v_i has a neighbour in $\{v_1, \dots, v_{i-1}\}$ and a neighbour in $\{v_{i+1}, \dots, v_n\}$.

Lemma 3 (Lempel et al. [15]). *A graph G is a 2-connected graph if, and only if, for every pair (s, t) of vertices, it admits an (s, t) -ordering.*

In fact, Lempel et al. established Lemma 3 only for simple graphs but it can be trivially extended to graphs since replacing all the parallel edges between two vertices by a unique edge does not change the connectivity.

Theorem 4. *Let G be a 2-connected subcubic graph. Then $c_3(G) \leq 3 \cdot 2^{n - \frac{n_3}{2}}$.*

Proof. If G is a cycle, then the result follows Proposition 1. Hence we may assume that G is not a cycle and thus has at least two vertices of degree 3, say s and t . By Lemma 3, there exists an (s, t) -ordering v_1, v_2, \dots, v_n of G . Orient all the edges of G according to this order, that is from the lower-indexed end-vertex towards its upper-indexed ones. Let us denote by D the obtained digraph. Observe that $d^+(v_1) = 3 = d^-(v_n)$ and $d^-(v_1) = 0 = d^+(v_n)$. Let A^+ (resp. A^- , A_2) be the set of vertices with outdegree 2 (resp. indegree 2) in D and A_2 be the set of vertices with degree 2 in G (and thus with indegree 1 and outdegree 1 in D). Clearly, (A_2, A^-, A^+) is a partition of $V(D) \setminus \{v_1, v_n\}$. Observe that $|A_2| = n - n_3$. Since $\sum_{v \in V(D)} d^+(v) = \sum_{v \in V(D)} d^-(v)$, we have $|A^+| = |A^-|$, and so $|A^+| = (n_3 - 2)/2$.

Now for $i = 1$ to $n - 1$, we enumerate the p_i partial 3-edge-colourings of the arcs with tail in $\{v_1, \dots, v_i\}$. For $i = 1$, there are 6 such colourings, since $d^+(v_1) = 3$.

Now, for each i , when we want to extend the partial colourings, two cases may arise.

- If $d_D^-(v_i) = 1$, then we need to colour one or two arcs, and one colour (the one of the arc entering v_i) is forbidden, so there are at most 2 possibilities. Hence $p_i \leq 2p_{i-1}$.
- If $d_D^-(v_i) = 2$, then we need to colour one arc, and at least two colours (the ones of the arcs entering v_i) is forbidden, so there is at most one possibility. Hence $p_i \leq p_{i-1}$.

At the end, all the edges of G are coloured, and a simple induction shows that $c_3(G) = p_{n-1} \leq 6 \times 2^{|A_2| + |A^+|} = 3 \times 2^{n - \frac{n_3}{2}}$. \square

In particular, for a cubic graph G , we obtain $c_3(G) \leq 3 \times 2^{n/2}$. We now extend this result to k -regular graphs.

Theorem 5. *Let G be a connected k -regular graph, with $k \geq 3$. Then $c_k(G) \leq k \cdot ((k-1)!)^{n/2}$.*

Proof. First, remark that if a k -regular graph G admits a k -edge-colouring, then every colour induces a perfect matching of G , and then n is even. Furthermore, observe that G is 2-connected. Indeed, assume that G has a cutvertex x and admits a k -edge-colouring c . As G has an even number of vertices, one of the connected components, say C , of $G - x$ has odd cardinality. A colour appearing on an edge between x and a connected component of $G - x$ different from C must form a perfect matching on C which is impossible. So, G is 2-connected.

Hence we can use the method of the proof of Theorem 4 and consider an (s, t) -ordering v_1, \dots, v_n of G and D the orientation of G obtained from this ordering (i.e. $v_i v_j \in A(D)$ if and only if $v_i v_j \in E(G)$ and $i < j$). The analysis made in the proof of Theorem 4 yields $c_k(G) \leq \prod_{x \in V(G)} (d^+(x)!)^i$. For $i = 1, \dots, k-1$, we define $A_i = \{x \in V(G) \setminus \{v_1, v_n\} : d^+(x) = i\}$. It is clear that $(A_i)_{1 \leq i \leq k-1}$ form a partition of $V(G) \setminus \{v_1, v_n\}$. If we denote $|A_i|$ by a_i , then $c_k(G) \leq P := k! \prod_{i=1}^{k-1} (i!)^{a_i}$. Moreover $S_1 := \sum_{i=1}^{k-1} a_i = n-2$ (by counting the number of vertices of G) and $S_2 := \sum_{i=1}^{k-1} i \cdot a_i = k(n-2)/2$ (by counting the number of arcs of $D - v_1$).

Let us now find the maximum value of P under the conditions $S_1 = n-2$ and $S_2 = k(n-2)/2$. If we can find $1 < p \leq q < k-1$ with $a_p \neq 0$ and $a_q \neq 0$ (or $a_p \geq 2$ if $p = q$), then we decrease a_p and a_q by one and increase a_{p-1} and a_{q+1} by one. Doing this, S_1 and S_2 are unchanged and P is multiplied by $\frac{q+1}{p} > 1$. We repeat this operation as many times as possible and stop when (a) for every $i = 2, \dots, k-2$, $a_i = 0$ or (b) there exists $j \in \{2, \dots, k-2\}$ such that for every $i = 2, \dots, k-2$ and $i \neq j$, $a_i = 0$ and $a_j = 1$. In case (b), S_1 gives $a_1 + 1 + a_{k-1} = n-2$ and S_2 is $a_1 + j + (k-1)a_{k-1} = k(n-2)/2$. Combining S_1 and S_2 , we obtain $2(k-2)a_{k-1} + 2(j-1) = (k-2)(n-2)$ and we conclude that $k-2$ divides $2(j-1)$ and so that $j = k/2$. Solving S_1 and S_2 we have in particular that $a_1 = a_{k-1}$ and that $2a_1 = n-3$ which is impossible, as n is even. Hence, we are in case (a), and solving S_1 and S_2 yields to $a_1 = a_k = (n-2)/2$. We conclude that $P \leq k!((k-1)!)^{n/2-1}$. \square

We turn now the proof of Theorem 5 into an algorithm to enumerate all the k -edge-colourings of a k -regular graph.

Corollary 6. *There is an algorithm to enumerate all the k -edge-colouring of a k -regular graph on n vertices in time $O^*((k-1)!)^{n/2}$ and polynomial space.*

Proof. Let G be a k -regular graph. First, we compute an (s, t) -ordering v_1, \dots, v_n of G , which can be done in polynomial time (see [5] and [6] for instance), and orient G accordingly to this ordering. Now, it is classical to enumerate all the permutations of a set of size p in time $O(p!)$ and linear space, in such way that, being given a permutation we compute in average constant time the next permutation in the enumeration (with the Steinhaus-Johnson-Trotter algorithm for instance, see [12]).

Using this and the odometer principle, it is now easy to enumerate all the edge colourings we want. In the enumeration of all the permutations of $\{1, \dots, k\}$, we take the first one and assign the corresponding colours to the arc

with tail x_1 . For any index i with $2 \leq i \leq k$, we assign to the arcs with tail x_i the first permutation in the enumeration of the permutations of the possible colours for these arcs (i.e. all the colours of $\{1, \dots, k\}$ minus the one of the arcs entering in x_i). Then, we have the first colouring, and we check if it is a proper edge colouring of G (in polynomial time). To have the following colouring, we take the next permutation on the colours possible on the arcs with tail x_{n-1} , and so on. Once all the possible permutations have been enumerated for these arcs, we take the next permutation on the colours possible on the arcs with tail x_{n-2} and re-enumerate the permutation of possible colours for the arcs with tail x_{n-1} , and so on, following the odometer principle. \square

The bound given by Theorem 5 is optimal on the class of k -regular graph. For all $k \geq 3$, and $n \geq 2$, $n = 2p$ even, the k -noodle necklace N_n^k is the k -regular graph obtained from a cycle on $2p$ vertices $(v_1, v_2, \dots, v_{2p}, v_1)$ by replacing all the edges $v_{2i-1}v_{2i}$, $1 \leq i \leq p$ by $k-1$ parallel edges.

Proposition 7. *Let $k \geq 3$ and $n \geq 2$,*

$$c_k(N_n^k) = k \cdot ((k-1)!)^{n/2}.$$

Proof. Observe that in every k -edge-colouring of N_n^k the edges which are not multiplied (i.e. $v_{2i}v_{2i+1}$) are coloured the same. There are k choices for such a colour. Once this colour is fixed, there are $(k-1)!$ choices for each set of $k-1$ parallel edges. Hence $c_k(N_n^k) = k \times ((k-1)!)^{n/2}$. \square

2.2 A more precise bound for cubic graphs

For simple cubic graph, we lower the bound on the number of 3-edge-colouring from $3 \cdot 2^{n/2}$ to $\frac{9}{4} \cdot 2^{n/2}$.

Lemma 8. *If G is a simple cubic graph, then $c_3(G) \leq \frac{9}{4} \cdot 2^{n/2}$.*

Proof. As in Theorem 4, let us consider an (s, t) -ordering v_1, v_2, \dots, v_n of the vertices and the acyclic digraph D obtained by orienting all the edges of G according to this ordering.

Let i be the smallest integer such that $d^-(v_i) = 2$. Since every vertex has an inneighbour, there exists j such that there are two internally-disjoint directed paths from v_j to v_i in D . In G , the union of these two paths forms a cycle C . By definition of i , all vertices of C but v_i have outdegree 2. So, if there is k such that $j < k < i$ and $v_k \notin V(C)$, then v_k has no outneighbour in C and the ordering $v_1, \dots, v_{k-1}, v_{k+1}, \dots, v_i, v_k, v_{i+1}, \dots, v_n$ is also an (s, t) -ordering. Repeating this operation as many times as necessary, we may obtain that all the vertices of C are consecutive in the ordering, that is $C = v_j v_{j+1} \dots v_i v_j$.

We enumerate the 3-edge-colourings of G in a similar way to the proof of Theorem 4, except that instead of examining the colour of arcs with tail in $\{v_j, \dots, v_i\}$ one after another, we look at C globally. If $j = 1$, then there are exactly $c_3(C)$ such 3-edge-colourings, because no arcs has head in C . If $j > 1$, then there are $c_3(C)/3$ such 3-edge-colourings, because one arc has head v_j and

we need the colours of the two arcs with tail v_j to have a colour distinct from it.

Recall that in D , v_1 has indegree 0, v_p has indegree 3, $\frac{n-2}{2}$ vertices have indegree 2 and $\frac{n-2}{2}$ vertices have indegree 1. If $j > 1$ (resp. $j = 1$), then there are $i - j$ (resp. $i - 2$) vertices of indegree 1 in C , so there are $\frac{n-2i+2j-2}{2}$ (resp. $\frac{n-2i+2}{2}$) vertices of indegree 1 in $V(G) \setminus V(C)$.

If $j = 1$, then we start by colouring C and then extend the colouring to G . Once C is coloured, there is at most one possibility to colour each arc with tail in C , so $c_3(G) \leq c_3(C) \cdot 2^{\frac{n-2i+2}{2}} = (c_3(C)/2^{i-1}) \cdot 2^{n/2}$. If $j > 1$, we colour the arcs with tail in $\{v_1, \dots, v_{j-1}\}$ as usual. Remark that, by the choice of C , there is exactly one of these arcs, denoted by e , which has head in C and more precisely e has tail v_j . Then, we consider all the edge colourings of C which agrees with the colour of e (i.e. $c(v_j v_{j+1}) \neq c(e)$ and $c(v_j v_i) \neq c(e)$). There is exactly $c_3(C)/3$ such colourings. Finally, we extend the edge colourings in all possible ways to G using the usual method. So, in this case, we obtain $c_3(G) \leq 6 \times c_3(C)/3 \times 2^{\frac{n-2i+2j-2}{2}} = (c_3(C)/2^{i-j}) \cdot 2^{n/2}$.

In all cases, we have to bound the value $\frac{c_3(C)}{2^{i-j}}$ for $1 \leq i < j$. Since G has no 2-cycles, C has length at least 3, and so $i - j \geq 2$. By Proposition 1, $c_3(G) = 2^{i-j+1} + 2$ if $i - j$ is odd, and $c_3(G) = 2^{i-j+1} - 2$ if $i - j$ is even. Easily one sees that the value $\frac{c_3(C)}{2^{i-j}}$ is maximized when $i - j = 3$ (i.e. C has length four), and so $\frac{c_3(C)}{2^{i-j}} \leq \frac{18}{8} = \frac{9}{4}$. Thus $c_3(G) \leq \frac{9}{4} \cdot 2^{n/2}$. \square

Theorem 5 for $k = 3$ states that a cubic graph G has at most $3 \cdot 2^{n/2}$ 3-edge-colourings. We shall now describe all graphs attaining this bound.

Let G be a cubic graph and $C = uvu$ be a 2-cycle in G . Then G/C is the graph obtained from $G - \{u, v\}$ by adding an edge between the neighbour of u distinct from v and the neighbour of v distinct from u .

Let Θ be the graph with two vertices joined by three edges. And let \mathcal{L} be the family of graphs defined recursively as follows:

- $\Theta \in \mathcal{L}$.
- if G has a 2-cycle C such that G/C is in \mathcal{L} , then G is in \mathcal{L} .

Remark that the k -noodle necklaces N_n^k (with n even) belongs to the family \mathcal{L} .

Theorem 9. *Let G be a cubic graph. If $G \in \mathcal{L}$, then $c_3(G) = 3 \cdot 2^{n/2}$. Otherwise $c_3(G) \leq \frac{9}{4} \cdot 2^{n/2}$.*

Proof. By induction on n , the result holding for simple graphs by Lemma 8 and for Θ because $c_3(\Theta) = 6$.

Assume that $n \geq 4$ and that G has a 2-cycle $C = uvu$. In any 3-edge-colouring of G , the edges not in C incident to u and v are coloured the same. Hence to each 3-edge-colouring c of G/C corresponds the two 3-edge-colourings of G that agrees with c on $G - \{u, v\}$. Hence $c_3(G) = 2c_3(G/C)$.

If G/C is in \mathcal{L} , then G is also in \mathcal{L} . Moreover, by induction hypothesis, $c_3(G/C) = 3 \cdot 2^{(n-2)/2}$. So $c_3(G) = 3 \cdot 2^{n/2}$.

If G/C is not in \mathcal{L} , then G is not in \mathcal{L} . Moreover, by induction hypothesis, $c_3(G/C) \leq \frac{9}{4} \cdot 2^{(n-2)/2}$. So $c_3(G) \leq \frac{9}{4} \cdot 2^{n/2}$. \square

We have no example of cubic simple graph admitting exactly $\frac{9}{4} \cdot 2^{n/2}$ 3-edge-colourings, and we believe that $\frac{9}{4}$ could be replaced by a lower constant in the statement of Theorem 9. In fact, we conjecture that the maximum number of 3-edge-colourings of cubic simple graphs of order n is attained by some special graphs that we now describe.

For all $n \geq 2$, $n = 2p$ even, the *hamster wheel* H_n is the cubic graph obtained from two cycles on p vertices $C_v = (v_1, v_2, \dots, v_p, v_1)$ and $C_w = (w_1, w_2, \dots, w_p, w_1)$ by adding the matching $M = \{v_i w_i : 1 \leq i \leq p\}$.

Proposition 10.

$$c_3(H_n) = \begin{cases} 2^{n/2} + 8, & \text{if } n/2 \text{ is even,} \\ 2^{n/2} - 2, & \text{if } n/2 \text{ is odd.} \end{cases}$$

Proof. Let ϕ is a 3-edge-colouring of C_v .

If the three colours appears on C_v , then there is a unique 3-edge-colouring of H_n extending ϕ . Indeed, to extend ϕ , the colours of the edges of M are forced. Since the three colours appears on C_v , there are two edges of M which are coloured differently. Without loss of generality, we may assume that these two edges are consecutive, that is there exists i such that they are $v_j w_j$ and $v_{j+1} w_{j+1}$. But then the colour of $w_j w_{j+1}$ must be equal to the one of $v_j v_{j+1}$. Then, from edge to edge along the cycle, one shows that for all i , the colour of $w_i w_{i+1}$ is the one of $v_i v_{i+1}$.

If only two colours appears on C_v , then there are two 3-edge-colourings of H_n extending ϕ . Indeed in this case, n is even and all the edges of M must be coloured by the colour not appearing on C_v . So, there are two possible 3-edge-colourings of C_w with the colours appearing on C_v .

Hence the number of 3-edge-colourings of G is equal to the number of 3-edge-colourings of C_v plus the number of 3-edge-colourings of C_v in which two colours appear. If $n/2$ is odd, this last number is 0, and if $n/2$ is even, this number is 6. So, by Proposition 1, $c_3(H_n) = 2^{n/2} - 2$ if $n/2$ is odd and $c_3(H_n) = 2^{n/2} + 8$ if $n/2$ is even. \square

For all $n \geq 2$, $n = 2p$ even, the *Möbius ladder* M_n is the cubic graph obtained from a cycle on n vertices $C = (v_1, v_2, \dots, v_n, v_1)$ by adding the matching $M = \{v_i v_{i+p} : 1 \leq i \leq p\}$ (indices are modulo n).

Two edges e and f of the cycle C are said to be *antipodal*, if there exists $1 \leq i \leq p$ such that $\{e, f\} = \{v_i v_{i+1}, v_{i+n} v_{i+p+1}\}$. A 3-edge-colouring c of M_n is said to be *antipodal* if $c(e) = c(f)$ for every pair (e, f) of antipodal edges.

Proposition 11. *Let c be a 3-edge-colouring of M_n . If c is not antipodal, then $n/2$ is odd and all the arcs $v_i v_{i+n/2+1}$ are coloured the same.*

Proof. Suppose that two antipodal edges are not coloured the same. Without loss of generality, $c(v_n v_1) = 2$ and $c(v_p v_{p+1}) = 3$, where $n = 2p$. Hence, we have $c(v_1 v_{p+1}) = 1$, $c(v_1 v_2) = 3$ and $c(v_{p+1} v_{p+2}) = 2$. And so on by induction, for all $1 \leq i \leq p$, $c(v_i v_{i+p}) = 1$ and $\{c(v_i v_{i+1}), c(v_{i+p} v_{i+p+1})\} = \{2, 3\}$. Hence, the edges of C are coloured alternately with 2 and 3. Since $c(v_{2p} v_1) = 2$ and $c(v_p v_{p+1}) = 3$, p must be odd. \square

Proposition 12.

$$c_3(M_n) = \begin{cases} 2^{n/2} + 2, & \text{if } n/2 \text{ is even,} \\ 2^{n/2} + 4, & \text{if } n/2 \text{ is odd.} \end{cases}$$

Proof. Clearly, there is a one-to-one mapping between the antipodal 3-edge-colourings of M_n and the 3-edge-colourings of $C_{n/2}$. Hence, by Proposition 11, if $n/2$ is even, then $c_3(M_n) = c_3(C_{n/2}) = 2^{n/2} + 2$ by Proposition 1.

If $n/2$ is odd, non-antipodal 3-edge-colourings are those such that all arcs $v_i v_{i+n/2+1}$ are coloured the same, by Proposition 11. There are 6 such edge colourings (three choices for the colour of the edges $v_i v_{i+n/2+1}$ and for each of these choices, two possible edge colourings of C). Hence $c_3(M_n) = c_3(C_{n/2}) + 6 = 2^{n/2} + 4$ by Proposition 1. \square

We think that H_n and M_n are the cubic graphs which admit the maximum number of 3-edge-colourings. Precisely, we raise the following conjecture.

Conjecture 13. Let G be a connected cubic simple graph on n vertices. If $n/2$ is even, then $c_3(G) \leq c_3(H_n)$ and if $n/2$ is odd, then $c_3(G) \leq c_3(M_n)$.

3 Total colouring

A *total colouring* of a graph G into k colours is a colouring of its vertices and edges such that two adjacent vertices receive different colours, two adjacent edges receive different colours and a vertex and an edge incident to it receive different colours. A total colouring with k colours is a *k-total-colouring*. For every graph G , let $c_k^T(G)$ be the number of k -total-colourings of G .

For each 4-edge-colouring c of a cubic graph G , there is at most one 4-total-colouring of G whose restriction to $E(G)$ equals c . Indeed, if the three edges adjacent to a vertex are coloured, the colour of the vertex is forced. Hence if G is cubic, we have that $c_4^T(G) \leq c_4(G)$.

By the method described in the previous section, one can show that if G is 2-connected, then $c_4(G) = O(2^{n/2} \cdot 6^{n/2})$, and so $c_4^T(G) = O(2^{n/2} \cdot 6^{n/2})$. We now obtain better upper bounds for c_4^T .

Theorem 14. Let G be a 2-connected subcubic graph. Then $c_4^T(G) \leq 3 \cdot 2^{2n-n_3/2}$.

Proof. Assume first that G is a cycle v_1, \dots, v_n . Let us total-colour it greedily starting from v_1 . There are 4 possible colours for v_1 , and then 3 possible colours

for v_1v_2 . Afterwards for every $i \geq 2$, there at most two possible colours for v_i (the ones distinct from the colours of v_{i-1} and $v_{i-1}v_i$) and then at most two possible colours for v_iv_{i+1} (the ones distinct from the colours of $v_{i-1}v_i$ and v_i). Hence $c_4^T(G) \leq 4 \times 3 \times 2^{2n-2} = 3 \cdot 2^{2n}$.

Assume now that G is not a cycle. Let s and t be two distinct vertices of order 3. Consider an (s, t) -ordering v_1, v_2, \dots, v_n of $V(G)$, which exists by Lemma 3, and the orientation D of G according to this ordering. Then $d^+(v_1) = 3 = d^-(v_n)$ and $d^-(v_1) = 0 = d^+(v_n)$. Let A^+ (resp. A^-) be the set of vertices of outdegree 2 (resp. indegree 2) in D and A_2 be the set of vertices of degree 2 in G . As in the proof of Theorem 4, we have $|A_2| = n - n_3$, and $|A^+| = |A^-| = (n_3 - 2)/2$.

Now for $i = 1$ to $n - 1$, we enumerate the p_i partial 4-total-colourings of vertices in $\{v_1, \dots, v_i\}$ and arcs with tail in $\{v_1, \dots, v_i\}$. For $i = 1$, there are $4! = 24$ such colourings, since v_1 and its three incident arcs must receive different colours.

For each $1 < i < n$, when we extend the partial total colourings. Two cases may arise.

- If $d_D^-(v_i) = 1$, then there are two choices to colour v_i and then two other choices to colour the (at most two) arcs leaving v_i . Hence $p_i \leq 4p_{i-1}$.
- If $d_D^-(v_i) = 2$, then there are at most two choices to colour v_i and then the colour of the arc leaving v_i is forced since three colours are forbidden by v_i and its two entering arcs. Hence $p_i \leq 2p_{i-1}$.

Finally, we need to colour v_n . Since its three entering arcs are coloured its colour is forced (or it is impossible to extend the colouring).

Hence an easy induction shows that $c_4^T(G) = p_{n-1} \leq 24 \cdot 4^{|A_2|+|A^+|} \cdot 2^{|A^-|} = 3 \cdot 2^{2n-n_3/2}$. \square

A *leaf* of a tree is a degree one vertex. A vertex of a tree which is not a leaf is called a *node*. A tree is *binary* if all its nodes have degree 3.

Proposition 15. *If T is a binary tree of order n , then $c_4^T(T) = 3 \cdot 2^{3n/2}$.*

Proof. By induction on n , the results holding easily when $n = 2$, that is when $T = K_2$.

Suppose now that T has more than two vertices. There is a node x which is adjacent to two leaves y_1 and y_2 . Consider the tree $T' = T - \{y_1, y_2\}$. By the induction hypothesis, $c_4^T(T') = 3 \cdot 2^{3(n-2)/2}$. Now each 4-total-colouring of T' may be extended into exactly eight 4-total-colourings of T' . Indeed the two colours of x and its incident edge in T' are forbidden for xy_1 and xy_2 , so there are two possibilities to extend the colouring to these edges, and then for each y_i , there are two possible colours available. Hence $c_4^T(T) = 8 \times c_4^T(T') = 3 \cdot 2^{3n/2}$. \square

Theorem 16. *Let G be a connected cubic graph. Then $c_4^T(G) \leq 3 \cdot 2^{3n/2}$.*

Proof. Let F be the subgraph induced by the cutedges of G . Then F is a forest. Consider a tree of F . It is binary, its leaves are in different non-trivial 2-connected components of G , and every node is a trivial 2-connected component of G .

A subgraph H of G is *full* if it is induced on G , connected and such that for every non-trivial 2-connected component C , $H \cap C$ is empty or is C itself and for every tree T of F , $H \cap T$ is empty, or is just one leaf of T or is T itself. Observe that a full subgraph has minimum degree at least 2.

We shall prove that for every full subgraph H , $c_4^T(H) \leq 3 \cdot 2^{2n(H) - n_3(H)/2}$. We proceed by induction on the number of 2-connected components of H . If H is 2-connected, then the result holds by Theorem 14.

Suppose now that H is not 2-connected. Then H contains a tree T of F . Let v_1, \dots, v_p be the leaves of T , e_i , $1 \leq i \leq p$ the edge incident to v_i in T and N the set of nodes of T . Then $H - N$ has p connected components H_1, \dots, H_p such that $v_i \in H_i$ for all $1 \leq i \leq p$. Furthermore, each H_i is a full subgraph of G .

Let c be a 4-total-colouring of T . It can be extended to H_i by any 4-total-colouring of H_i such that v_i is coloured $c(v_i)$ and the two edges incident to v_i in H_i are coloured in $\{1, 2, 3, 4\} \setminus \{c(v_i), c(e_i)\}$. There are $\frac{1}{12}c_4^T(H_i)$ such colourings ($\frac{1}{4}$ to choose $c(v_i)$ times $\frac{1}{3}$ to choose $c(e_i)$). Hence each 4-total-colouring of T can be extended into $\prod_{i=1}^p \frac{1}{12}c_4^T(H_i)$ 4-total-colourings of H and so

$$c_4^T(H) = c_4^T(T) \times \prod_{i=1}^p \frac{1}{12}c_4^T(H_i).$$

Now by Proposition 15, T has $3 \cdot 2^{3n(T)/2}$ 4-total-colourings, and since H_i is full $c_4^T(H_i) \leq 3 \cdot 2^{2n(H_i) - n_3(H_i)/2}$ by the induction hypothesis. Moreover, $n_3(H) = n(T) + \sum_{i=1}^p n_3(H_i)$ and $n(H) = n(T) + \sum_{i=1}^p n(H_i) - p$. Hence $c_4^T(H) \leq 3 \cdot 2^{2n(H) - n_3(H)/2}$. \square

As previously for the edge colourings of graphs, we derive from Theorem 16 an algorithm to enumerate all the 4-total-colourings of a cubic graph. The proof is similar to the one of Corollary 6.

Corollary 17. *There is an algorithm to enumerate all the 4-total-colourings of a cubic graph on n vertices in time $O^*(2^{3n/2})$ and polynomial space.*

The bound of Theorem 16 is seemingly not tight. Indeed, in Theorem 14, the equation $p_i \leq 2p_{i-1}$ when $d_D^-(v_i)$ often overestimate p_i , because there are two choices to colour v_i only if the two colours appearing on its two entering arcs are the same two as the ones assigned to the tails of these arcs. If not the colour of v_i is forced or v_i cannot be coloured.

Problem 18. What is $c_4^T(n)$, the maximum of $c_4^T(G)$ over all graphs of order n ?

We shall now give a lower bound on $c_4^T(n)$. A binary tree is *nice* if its set of leaves may be partitionned into pairs of *twins*, i.e. leaves at distance 2. Clearly,

every nice binary tree T has an even number of leaves and thus $n(T) \equiv 2 \pmod{4}$. Moreover if $n(T) = 4p + 2$, then T has $2p$ nodes and $p + 1$ pairs of twins. A *noodle tree* is a cubic graph obtained from a nice binary tree by adding two parallel edges between each pair of twins.

Proposition 19. *Let p be a positive integer and $n = 4p + 2$. If G is a noodle tree G of order n , then $c_4^T(G) = \frac{3}{\sqrt{2}} \cdot 2^{5n/4}$.*

Proof. Let X_1, \dots, X_{p+1} be the pairs of twins of G , and let T be the binary tree $G - \bigcup_{i=1}^{p+1} X_i$. Let us label the leaves of T , y_1, \dots, y_{p+1} such that for all $1 \leq i \leq p + 1$, y_i is adjacent to the two vertices of X_i in G .

Every 4-total-colouring of T , may extended in exactly 4 ways to each pair of twins $X_i = \{x_i, x'_i\}$ and their incident edges. Indeed, without loss of generality we may assume that y_i is coloured 1 and its incident edge in T is coloured 2. Then the edges $y_i x_i$ and $y_i x'_i$ must be coloured in $\{3, 4\}$, which can be done in two possible ways. For each of these possibilities, the parallel edges between x_i and x'_i must be coloured in $\{1, 2\}$, which again can be done in two possible ways. Finally, we must colour x_i (resp. x'_i) with the colour of $y_i x'_i$ (resp. $y_i x_i$).

Hence $c_4(G) = 4^{p+1} \times c_4(T)$, and so by Proposition 15, $c_4(G) = 3 \cdot 2^{5p+2}$. \square

Acknowledgement

The authors would like to thank their children for suggesting them some graph names.

References

- [1] R. Beigel and D. Eppstein. 3-coloring in time $O(1.3289^n)$. *Journal of Algorithms* 54:168–204, 2005.
- [2] A. Björklund and T. Husfeldt. Inclusion-exclusion algorithms for counting set partitions. In *Proceedings of the 47th Annual IEEE Symposium on Foundations of Computer Science (FOCS 2006)*, pp. 575–582. IEEE, Los Alamitos, 2006.
- [3] A. Björklund, T. Husfeldt, P. Kaski and M. Koivisto. Narrow sieves for parameterized paths and packings. arXiv:1007.1161v1, 2010.
- [4] J. A. Bondy and U. S. R. Murty. *Graph theory*. Graduate Texts in Mathematics 244, Springer, 2008 (2nd ed.).
- [5] S. Even and R. E. Tarjan. Computing an st-numbering. *Theoretical Computer Science*, 2(3):339344, 1976.
- [6] S. Even and R. E. Tarjan. Corrigendum: Computing an st-numbering. *Theoretical Computer Science*, 4(1):123, 1977.

-
- [7] F. V. Fomin, S. Gaspers, and S. Saurabh. Improved exact algorithms for counting 3- and 4-colorings. In *Lin, G. (ed.) COCOON 2007, Lecture Notes in Computer Science* 4598:65–74, 2007.
 - [8] M. R. Garey and D. S. Johnson, *Computers and intractability. A guide to the theory of NP-completeness*. A Series of Books in the Mathematical Sciences. W. H. Freeman and Co., San Francisco, Calif., 1979.
 - [9] P. A. Golovach, D. Kratsch, and J. F. Couturier. Colorings with few colors: counting, enumeration and combinatorial bounds. In *Proceedings of WG 2010, Lecture Notes in Computer Science* 6410:39–50, 2010.
 - [10] I. Holyer. The NP-completeness of edge-colouring. *SIAM J. Computing* 2:225–231, 1981.
 - [11] T. R. Jensen and B. Toft. *Graph coloring problems*. Wiley-Interscience Series in Discrete Mathematics and Optimization. John Wiley & Sons, Inc., New-York, 1995.
 - [12] D. E. Knuth. *The art of Computer Programming, volume 4A: Combinatorial Algorithms, Part 1*. Addison-Wesley Professional, 2011.
 - [13] M. Koivisto. An $O(2^n)$ algorithm for graph coloring and other partitioning problems via inclusion-exclusion. In *Proceedings of the 47th Annual IEEE Symposium on Foundations of Computer Science (FOCS 2006)*, pp. 583–590. IEEE, Los Alamitos, 2006.
 - [14] L. Kowalik. Improved edge-coloring with three colors. *Theoret. Comp. Sci.* 410:3733–3742, 2009.
 - [15] A. Lempel, S. Even, and I. Cederbaum. An algorithm for planarity testing of graphs. In *P. Rosenstiehl, editor, Proceedings of the International Symposium on the Theory of Graphs (Rome, July 1966)*: 215–232. Gordon and Breach, 1967.
 - [16] A. Sánchez-Arroyo. Determining the total colouring number is NP-hard. *Discrete Math.* 78(3):315–319, 1989.



Centre de recherche INRIA Sophia Antipolis – Méditerranée
2004, route des Lucioles - BP 93 - 06902 Sophia Antipolis Cedex (France)

Centre de recherche INRIA Bordeaux – Sud Ouest : Domaine Universitaire - 351, cours de la Libération - 33405 Talence Cedex
Centre de recherche INRIA Grenoble – Rhône-Alpes : 655, avenue de l'Europe - 38334 Montbonnot Saint-Ismier
Centre de recherche INRIA Lille – Nord Europe : Parc Scientifique de la Haute Borne - 40, avenue Halley - 59650 Villeneuve d'Ascq
Centre de recherche INRIA Nancy – Grand Est : LORIA, Technopôle de Nancy-Brabois - Campus scientifique
615, rue du Jardin Botanique - BP 101 - 54602 Villers-lès-Nancy Cedex
Centre de recherche INRIA Paris – Rocquencourt : Domaine de Voluceau - Rocquencourt - BP 105 - 78153 Le Chesnay Cedex
Centre de recherche INRIA Rennes – Bretagne Atlantique : IRISA, Campus universitaire de Beaulieu - 35042 Rennes Cedex
Centre de recherche INRIA Saclay – Île-de-France : Parc Orsay Université - ZAC des Vignes : 4, rue Jacques Monod - 91893 Orsay Cedex

Éditeur
INRIA - Domaine de Voluceau - Rocquencourt, BP 105 - 78153 Le Chesnay Cedex (France)
<http://www.inria.fr>
ISSN 0249-6399